

Introduction to Cryptology

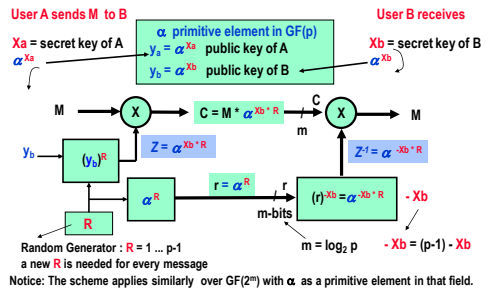
Tutorial-09

EIGamal Public-Key Systems over GF(p) & GF(2^m)

15.05.2023, v52

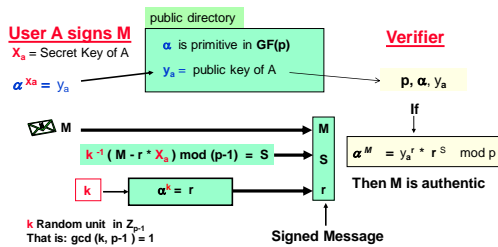
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EIGamal Secrecy-System (1985)



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EIGamal Signature Scheme



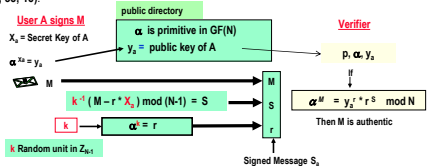
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EIGamal Secrecy-System Over GF (p)

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Problem 9-1: El-Gamal crypto system is set up using the prime number $N = 2 * 19 * 11 + 1 = 419$ generated by applying Pocklington's Theorem, where 19 and 11 are two primes.

1. Prove that N is a prime according to Pocklington's Theorem.
2. Compute the probability that a randomly selected element is primitive in GF(419).
3. Find a primitive element in GF(419) and use it as a public element in El-Gamal public key system.
4. User A encrypts the message $M = 21$ and send it to user B who has the secret key $X_b = 80$ by using the random number $R = 13$. Compute B's public key Y_b and the encrypted message C_b and r .
5. Decrypt the cryptogram C_b on the receiver side B showing all therefore necessary computations.
6. Let user A having the secret key $X_a = 133$ compute his Signature S_a according to El-Gamal signature scheme shown below for the same message $M = 21$. Select one adequate k from the following list ($k = 22, 38, 15$).



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Solution 9-1:

1. $N = R * F + 1 = 2 * 19 * 11 + 1 = 419$, $F = 19 * 11$ and $R = 2$. $p_1 = 19$, $p_2 = 11$ Is 419 a prime?

Proof: We select $a = 2$

1. $\text{gcd}(a^{(N-1)/p_1} - 1, N) = \text{gcd}(2^{22} - 1, 419) = 1$ is true
 $\text{gcd}(a^{(N-1)/p_2} - 1, N) = \text{gcd}(2^{38} - 1, 419) = 1$ is true

2. $a^{N-1} = 1 \text{ (mod } N) \Leftrightarrow 2^{418} = 1 \text{ (mod } 419)$ is true

3. $F > \sqrt{N}$

$11 * 19 > \sqrt{419} = 20.46 \Rightarrow 209 > 20.46$ is true
 As all conditions 1, 2 and 3 are true $\Rightarrow 419$ is prime

2. # of all non-zero elements : $419 - 1 = 418$
 # of primitive elements: $\phi(418) = \phi(2 * 11 * 19) = (2-1) * (11-1) * (19-1) = 180$
 $P(\text{element-primitive}) = (180 / 418) * 100 = 43.06\%$
3. Possible orders are the divisors of $419 - 1 = 418$
 These are: 1, 2, 11, 19, 22, 38, 209 and 418

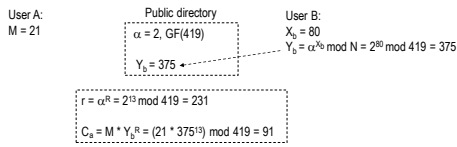
Checking if the element 2 is primitive

- $2^1 \text{ mod } 419 \neq 1$, $2^{22} \text{ mod } 419 \neq 1$
 $2^2 \text{ mod } 419 \neq 1$, $2^{38} \text{ mod } 419 \neq 1$
 $2^{11} \text{ mod } 419 \neq 1$, $2^{209} \text{ mod } 419 \neq 1 \Rightarrow \text{Ord}(2) = 418 \Rightarrow 2$ is primitive element
 $2^{19} \text{ mod } 419 \neq 1$

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Solution 9-1 cont.:

4. Encryption:



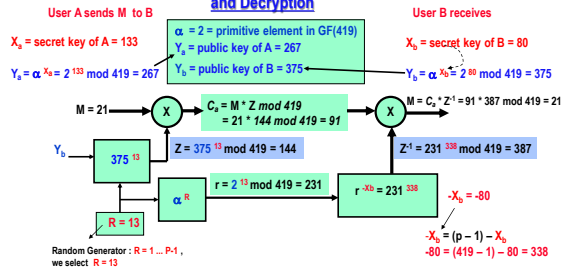
5. Decryption:

$$Z^1 = r^{-X_b} = 231^{-80} \pmod{419} = 231^{80-418} \pmod{419} = 387$$

$$M = C_b * Z^1 = 91 * 387 \pmod{419} = 21$$

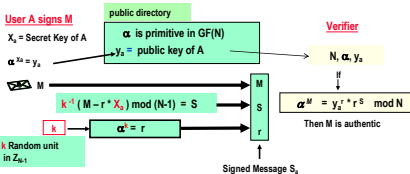
Solution 9-1 cont.:

Overview of Encryption and Decryption



Solution 9-1 cont.:

6.



k has to be invertible mod N-1, N-1 = 418
 $\text{gcd}[k, N-1] = 1 \Rightarrow$ select 15 as $\text{gcd}(418, 15) = 1$
 $k = 15, k^{-1} = 15^{-1} = -195 \pmod{418} = 223$ (see table below)
 $r = \alpha^k = 2^{15} \pmod{419} = 86$

$$S = k^{-1} (M - r * X_a) \pmod{N-1}$$

$$= 15^{-1} (21 - 86 * 133) \pmod{418}$$

$$= 223 (21 - 11438) \pmod{418}$$

$$= -2545991 \pmod{418}$$

$$S = 47$$

m	10	a1	a2	b1	b2	a	r	INVERSE VALUE	GG
418	15	86	0	1	27	13			
15	13	0	1	-27	1	2			
13	2	1	-13	-27	-28	6	1		
2	-1	-13	7	28	-195	2	0	INVERSE	-195 GG+

Problem 9-2: El-Gamal crypto system is set up using the prime number N = 29.

- Find a primitive element in GF(29) and use it as a public element in El-Gamal public key system.
- User A has the secret key $X_a = 7$ and User B has the secret key $X_b = 4$. Compute A's public key Y_a and B's public key Y_b .
- User A encrypts the message $M = 17$ and send it to user B. Compute the encrypted message C_b and r.
 - Use the random number $R = 21$.
 - Use the random number $R = 25$.
 - Which random number is better? Why?
- Decrypt the cryptogram C_b on the receiver side B showing all necessary computations therefore. Use the better random number R.

Solution 9-2:

- Possible orders are the divisors of $29 - 1 = 28$
 These are: 1, 2, 4, 7, 14 and 28

Checking if the element 3 is primitive

$$3^1 \pmod{29} \neq 1$$

$$3^2 \pmod{29} \neq 1$$

$$3^4 \pmod{29} \neq 1$$

$$3^7 \pmod{29} \neq 1$$

$$3^{14} \pmod{29} \neq 1 \Rightarrow \text{Ord}(3) = 28 \Rightarrow 3 \text{ is primitive element}$$

Solution 9-2 cont.:

2. User A: $X_a = 7, Y_a = \alpha^{X_a} \pmod{N} = 3^7 \pmod{29} = 12$
 User B: $X_b = 4, Y_b = \alpha^{X_b} \pmod{N} = 3^4 \pmod{29} = 23$

3.a. User A: $X_a = 7, M = 17, R = 21$
 Public directory: $\alpha = 2, \text{GF}(29), Y_a = 12, Y_b = 23$
 $r = \alpha^R = 3^{21} \pmod{29} = 17$
 $C_b = M * Y_b^r = (17 * 23^{21}) \pmod{29} = 17$

3.b. User A: $X_a = 7, M = 17, R = 25$
 Public directory: $\alpha = 2, \text{GF}(29), Y_a = 12, Y_b = 23$
 $r = \alpha^R = 3^{25} \pmod{29} = 14$
 $C_b = M * Y_b^r = (17 * 23^{25}) \pmod{29} = 21$

Solution 9-2 cont.:

3.c. The random number $R = 25$ is better, because the random number $R = 21$ has the consequence that $C_b = M$. So there is no encryption, when using $R = 21$.

5. Decryption:

$$Z^1 = r^{-X_b} = 14^{-4} \pmod{29} = 14^{4-28} \pmod{29} = 16$$

$$M = C_b * Z^1 = 21 * 16 \pmod{29} = 17$$

Solution 9-2 cont.:

Overview of Encryption and Decryption

User A sends M to B

$X_A = \text{secret key of A} = 7$
 $Y_B = \alpha^{X_A} = 3^7 \bmod 29 = 12$

$\alpha = 3 = \text{primitive element in GF}(29)$
 $Y_A = \text{public key of A} = 12$
 $Y_B = \text{public key of B} = 23$

User B receives

$X_B = \text{secret key of B} = 4$
 $Y_B = \alpha^{X_B} = 3^4 \bmod 29 = 23$

$M = 17$

$C_B = M * Z \bmod 29 = 17 * 20 \bmod 29 = 21$

$Z = 23^{25} \bmod 29 = 20$

$M = C_B * Z^{-1} = 21 * 16 \bmod 29 = 17$

$Z^{-1} = 14^{24} \bmod 29 = 16$

$r = 3^{25} \bmod 29 = 14$

$r^{-X_B} = 14^{24}$

$-X_B = -4$
 $-4 = (29 - 1) - X_B = 24$

Random Generator: $R = 1 \dots P-1$, we select $R = 25$ (aus 3.b.)

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Problem 9-3: El-Gamal crypto system is set up using the prime number $N = 2 * 131 + 1 = 263$ generated by applying Pocklington's Theorem, where 131 is a prime.

- Prove that N is a prime according to Pocklington's Theorem.
- Compute the probability that a randomly selected element is primitive in $GF(263)$.
- Find a primitive element in $GF(263)$ and use it as a public element in El-Gamal public key system.
- User A encrypts the message $M = 35$ and send it to User B who has the secret key $X_B = 113$ by using the random number $R = 22$. Compute B's public key Y_B and the encrypted message C_B and r.
- Decrypt the cryptogram C_B on the receiver side B showing all necessary computations therefore.
- Let user A having the secret key $X_A = 40$ compute his Signature S_A according to El-Gamal signature scheme shown below for the same message $M = 35$. Choose $k = 121$.

User A signs M

$X_A = \text{Secret Key of A}$
 $\alpha = 7$
 $Y_A = \text{public key of A}$

Verifier

P, α, Y_A

M

$k^{-1} * (M - r * X_A) \bmod (N-1) = S$

$\alpha^M = Y_A^r * r^S \bmod N$

Then M is authentic

Signed Message S_A

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Solution 9-3:

- $N = R * F + 1 = 2 * 131 + 1 = 263$, $F = p = 131$ and $R = 2$. Is 263 a prime?

Proof: We select $a = 11$

 - $\gcd(a^{N-1} - 1, N) = \gcd(11^{262} - 1, 263) = 1$ is true
 - $a^{N-1} = 1 \pmod N \Leftrightarrow 11^{262} = 1 \pmod{263}$ is true
 - $F > \sqrt{N}$
 $131 > \sqrt{263} = 16.22$ is true
 As all conditions 1, 2 and 3 are true \Rightarrow 263 is prime
- # of all non-zero elements: $263 - 1 = 262$
 # of primitive elements: $\phi(262) = \phi(2 * 131) = (2 - 1) * (131 - 1) = 130$
 $P(\text{element-primitive}) = (130 / 262) * 100 = 49.62\%$
- Possible orders are the divisors of $263 - 1 = 262$, these are: 1, 2, 131 and 262

Checking if the element 12 is primitive
 $12^1 \bmod 263 \neq 1$, $12^2 \bmod 263 \neq 1$
 $12^{131} \bmod 263 = 1 \Rightarrow \text{Ord}(12) = 131 \Rightarrow 12$ is not a primitive element

Checking if the element 7 is primitive
 $7^1 \bmod 263 \neq 1$, $7^2 \bmod 263 \neq 1$
 $7^{131} \bmod 263 \neq 1 \Rightarrow \text{Ord}(7) = 262 \Rightarrow 7$ is a primitive element

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Solution 9-3 cont.:

- Encryption:

User A: $M = 35$

Public directory: $\alpha = 7, GF(263)$
 $Y_B = 236$

User B: $X_B = 113$
 $Y_B = \alpha^{X_B} \bmod N = 7^{113} \bmod 263 = 236$

$r = \alpha^R = 7^{22} \bmod 263 = 11$
 $C_B = M * Y_B^r = (35 * 236^{22}) \bmod 263 = 16$
- Decryption:

$Z^{-1} = r^{-X_B} = 11^{-113} \bmod 263 = 11^{113-262} \bmod 263 = 183$
 $M = C_B * Z^{-1} = 16 * 183 \bmod 263 = 35$

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Solution 9-3 cont.:

Overview of Encryption and Decryption

User A sends M to B

$X_A = \text{secret key of A} = 40$
 $Y_B = \alpha^{X_A} = 7^{40} \bmod 263 = 166$

$\alpha = 7 = \text{primitive element in GF}(263)$
 $Y_A = \text{public key of A} = 166$
 $Y_B = \text{public key of B} = 236$

User B receives

$X_B = \text{secret key of B} = 113$
 $Y_B = \alpha^{X_B} = 7^{113} \bmod 263 = 236$

$M = 35$

$C_B = M * Z \bmod 263 = 35 * 23 \bmod 263 = 16$

$Z = 236^{22} \bmod 263 = 23$

$M = C_B * Z^{-1} = 16 * 183 \bmod 263 = 35$

$Z^{-1} = 11^{149} \bmod 263 = 183$

$r = 7^{22} \bmod 263 = 11$

$r^{-X_B} = 11^{149}$

$-X_B = -113$
 $-113 = (263 - 1) - X_B = 149$

Random Generator: $R = 1 \dots P-1$, we select $R = 22$

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Solution 9-3 cont.:

6.

User A signs M

$X_A = \text{Secret Key of A}$
 $\alpha = 7$
 $Y_A = \text{public key of A}$

Verifier

N, α, Y_A

M

$k^{-1} * (M - r * X_A) \bmod (N-1) = S$

$\alpha^M = Y_A^r * r^S \bmod N$

Then M is authentic

Signed Message S_A

k has to be invertible mod $N-1$, $N-1 = 262$
 $\gcd[k, N-1] = 1 \Rightarrow \gcd(121, 262) = 1$
 $k = 121, k^{-1} = 121^{-1} = 13$ (see table below)
 $r = \alpha^k = 7^{121} \bmod 263 = 85$

$S = k^{-1} * (M - r * X_A) \bmod (N-1)$
 $= 121^{-1} * (35 - 85 * 113) \bmod 262$
 $= 13 * (35 - 9605) \bmod 262$
 $= -124410 \bmod 262$
 $= 540$

m	v	a1	a2	b1	b2	a1 r	INVERSE VALUE	v2
418	151	1	0	0	1	127	13	
15	113	0	1	1	0	121	13	
13	2	1	1	-273	23	0	1	
2	1	1	7	28	-193	2	0	INVERSE: -193

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EIGamal Secrecy-System Over GF(2^m)

Problem 9-4: El-Gamal crypto system is set up using GF(2⁶), which is generated by the irreducible polynomial P(x) = x⁶ + x⁵ + x⁴ + x² + 1 = 1110101.

1. Check if you can take $\alpha = 0011$ as a primitive element.
2. User A encrypts the message $M = \alpha^3$ and send it to user B who has the secret key $X_b = 10$ by using the random number $R = 43$. Compute B's public key Y_b and the encrypted message C_a and r .
3. Decrypt the cryptogram C_a on the receiver side B showing all necessary computations therefore.

Note: For the selected P(x), e = 21 this mean ord(x) = 21 (from the table list of all irreducible polynomials over GF(2))

Helping computations :

$$\begin{aligned} x^6 &= x^5 + x^4 + x^2 + 1 & x^{11} &= x^6 + x^2 \\ x^7 &= x^6 + x^5 + x^2 + x + 1 & x^{12} &= x^6 + x^4 + x^3 + x^2 + 1 \\ x^8 &= x^6 + x^4 + x^3 + x^2 + x & x^{16} &= x^3 + x^2 + 1 \\ x^9 &= x^3 + 1 & x^{21} &= 1 \\ x^{10} &= x^4 + x \end{aligned}$$

Solution 9-4:

1. Possible orders are the divisors of 2⁶ - 1 = 63, these are: 1, 3, 7, 9, 21 and 63

Checking if the element $\alpha = 000011 = x + 1$ is primitive

$$\alpha^3 = (x + 1)^3 = (x + 1)^2(x + 1) = (x^2 + 1)(x + 1)$$

$$= x^2 + x^2 + x + 1 \neq 1$$

$$\alpha^7 = (x + 1)^7 = (x + 1)^4(x + 1)^3 = (x + 1)^4(x^3 + x^2 + x + 1)$$

$$= x^4 + x^2 + 1 \neq 1$$

$$\alpha^9 = (x + 1)^9 = (x + 1)^7(x + 1)^2$$

$$= x^2 + x^4 + x^2 \neq 1$$

$$\alpha^{21} = \alpha^{12} \cdot \alpha^9 = [\alpha^{6^2} \cdot \alpha^9 = (x^6)^2(x^6 + x^4 + x^2) = x^{10}(x^6 + x^4 + x^2) = (x^4 + x)(x^6 + x^4 + x^2)$$

$$= x^4 + x^2 + x^2 + x + 1 \neq 1$$

⇒ Ord(α) = 63 ⇒ α is a primitive element

Solution 9-4 cont.:

2. Encryption:

User A:
M = α^3

Public directory
 $\alpha = x + 1, GF(2^6)$

$$Y_b = \alpha^{10}$$

User B:
 $X_b = 10$

$$Y_b = \alpha^{10} = \alpha^{10}$$

Cryptogram:

$$C_a = M \cdot Y_b^R = \alpha^3 \cdot (\alpha^{10})^{43}$$

$$= \alpha^3 \cdot \alpha^{430 \text{ mod } 63} = \alpha^3 \cdot \alpha^{32} = \alpha^{35}$$

$$r = \alpha^R = \alpha^{43}$$

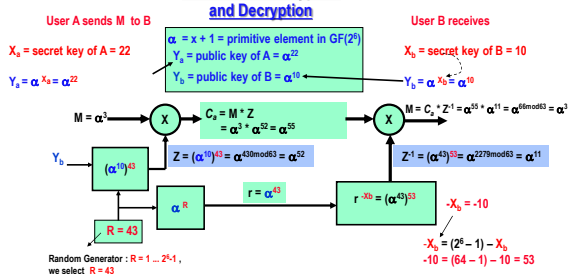
3. Decryption:

$$Z^1 = r^{-X_b} = (\alpha^{43})^{-10 \text{ mod } 63} = (\alpha^{43})^{-10+63} = (\alpha^{43})^{53} = \alpha^{2279 \text{ mod } 63} = \alpha^{11}$$

$$M = C_a \cdot Z^1 = \alpha^{35} \cdot \alpha^{11} = \alpha^{46 \text{ mod } 63} = \alpha^3$$

Solution 9-4 cont.:

Overview of Encryption and Decryption



Problem 9-5: El-Gamal crypto system is set up using GF(2⁶), which is generated by the irreducible polynomial P(x) = x⁶ + x + 1 = 10011.

1. Compute the exponents of the element $\beta = x = 000010$ as $x^i \text{ mod } P(x)$ for $i = 1$ to 15.
2. Which multiplicative orders are possible in GF(2⁶)?
3. Check if you can take $\alpha = 1011$ as a primitive element.
4. User A has the secret key $X_a = 7$ and User B has the secret key $X_b = 12$. Compute the public keys Y_a and Y_b .
5. User A encrypts the message $M = 0101$ and send it to user B by using the random number $R = 13$. Compute the encrypted message C_a and r .
6. Decrypt the cryptogram C_a on the receiver side B showing all necessary computations therefore.

Note: For the selected P(x), e = 15 this mean ord(x) = 15 (from the table list of all irreducible polynomials over GF(2))

Solution 9-5:

- If $P(x) = x^4 + x + 1$ is the modulus then $x^4 + x + 1 = 0$, thus $x^4 = x + 1$. The exponents of x in $GF(2^4)$ are:

$x = x$	0010
$x^2 = x^2$	0100
$x^3 = x^3$	1000
$x^4 = x + 1$	0011
$x^5 = x \cdot x^4 = x^2 + x$	0110
$x^6 = x \cdot (x^2 + x) = x^3 + x^2$	1100
$x^7 = x \cdot (x^3 + x^2) = (x^4 + x^2) = x^2 + x + 1$	1011
$x^8 = x^4 + x^2 + x = x + 1 + x^2 + x = x^2 + 1$	0101
$x^9 = x^3 + x$	1010
$x^{10} = x^4 + x^2 = x^2 + x + 1$	0111
$x^{11} = x^3 + x^2 + x$	1110
$x^{12} = x^4 + x^3 + x^2 = x^3 + x^2 + x + 1$	1111
$x^{13} = x^4 + x^3 + x^2 + x = x^3 + x^2 + 1$	1101
$x^{14} = x^4 + x^3 + x = x + 1 + x^3 + x = x^3 + 1$	1001
$x^{15} = x^4 + x = x + 1 + x = 1$	0001

 $\Rightarrow \text{ord}(x) = 15$
- Possible orders are the divisors of $2^4 - 1 = 15$, these are: 1, 3, 5 and 15

Solution 9-5 cont.:

- Checking if the element $\alpha = 1011 = x^7$ is primitive, $\text{ord}(x) = 15$ (aus 2.)

$$\text{ord}(\alpha) = \text{ord}(x^7) = \text{ord}(x) / \text{gcd}(\text{ord}(x), 7) = 15 / \text{gcd}(15, 7) = 15 / 1 = 15$$

$$\Rightarrow \text{ord}(\alpha) = 15 \Rightarrow \alpha \text{ is a primitive element}$$
- | | | | |
|---|-----------------------|--|------------|
| User A: | $\alpha = 1011 = x^7$ | User B: | $X_b = 12$ |
| $X_a = 7$ | | $Y_b = \alpha^{X_b} = (x^7)^{12} = x^{84} \text{ mod } 15 = x^9 = x^3 + x$ | |
| $Y_a = \alpha^{X_a} = (x^7)^7 = x^{49} \text{ mod } 15 = x^4 = x + 1$ | | $Y_b = \alpha^{X_b} = (x^7)^{12} = x^{84} \text{ mod } 15 = x^9 = x^3 + x$ | |
| | $= 0011$ | | $= 1010$ |

Solution 9-5 cont.:

- Encryption:

User A:	$M = 0101 = x^2 + 1 = x^8$	Public directory	User B:	$X_b = 12$
		$\alpha = x^7, GF(2^4)$		$Y_b = x^9$ (aus 4.)
		$Y_b = x^9$		

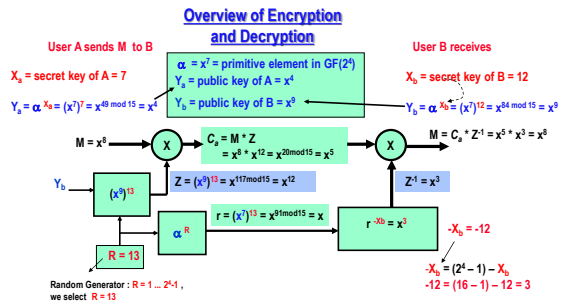
$$r = \alpha^R = (x^7)^{13} = x^{91} \text{ mod } 15 = x$$

$$C_a = M \cdot Y_b \cdot R = x^8 \cdot x^9 \cdot (x^9)^{13} = x^8 \cdot x^9 \cdot x^{117} \text{ mod } 15 = x^8 \cdot x^9 \cdot x^{12} = x^{29} \text{ mod } 15 = x^3 = x^2 + x = 0110$$
- Decryption:

$$Z^1 = r^{-X_b} = x^{-12} \text{ mod } 15 = x^{-12+15} = x^3$$

$$M = C_a \cdot Z^1 = x^8 \cdot x^3 = x^{11} = x^2 + 1 = 0101$$

Solution 9-5 cont.:



Problem 9-6: El-Gamal crypto system is set up using $GF(2^6)$, which is generated by the irreducible polynomial $P(x) = x^6 + x^2 + 1 = 1001001$.

- Compute the exponents of the element $\delta = x = 000010$ as $x^i \text{ mod } P(x)$ for $i = 1$ to 10.
 - Which multiplicative orders are possible in $GF(2^6)$?
 - Compute the probability that a randomly selected element is primitive in $GF(2^6)$.
 - Check if you can take $\alpha = x + 1$ as a primitive element.
 - User A has the secret key $X_a = 22$ and User B has the secret key $X_b = 10$. Compute the public keys Y_a and Y_b .
 - User A encrypts the message $M = 100100 = x^4 + x^2$ and send it to user B by using the random number $R = 20$. Compute the encrypted message C_a and r .
 - Decrypt the cryptogram C_a on the receiver side B showing all necessary computations therefore.
- Note:** For the selected $P(x)$, $e = 9$ this mean $\text{ord}(x) = 9$ (from the table list of all irreducible polynomials over $GF(2)$)

Solution 9-6:

- If $P(x) = x^6 + x^2 + 1$ is the modulus then $x^6 + x^2 + 1 = 0$, thus $x^6 = x^2 + 1$. The exponents of x in $GF(2^6)$ are:

$x = x$	000010
$x^2 = x^2$	000100
$x^3 = x^3$	001000
$x^4 = x^4$	010000
$x^5 = x^5$	100000
$x^6 = x^2 + 1$	001001
$x^7 = x \cdot (x^2 + 1) = x^3 + x$	010010
$x^8 = x^3 + x^2$	100100
$x^9 = x^5 + x^3 + x^2 + 1 + x^3 = 1$	000001
$x^{10} = x$	000010

 $\Rightarrow \text{ord}(x) = 9 \Rightarrow x$ is not primitive
- Possible orders are the divisors of $2^6 - 1 = 63$, these are: 1, 3, 7, 9, 21 and 63
- # of all non-zero elements: $2^6 - 1 = 63$
 # of primitive elements: $\phi(63) = \phi(3^2 \cdot 7) = 63 \cdot (1 - 1/3) \cdot (1 - 1/7) = 36$
 $P(\text{element} = \text{primitive}) = (36 / 63) \cdot 100 = 57.14\%$

Solution 9-7 cont.:

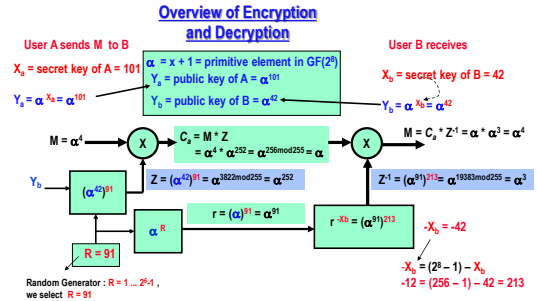
4. $\alpha = x + 1$
 User A: $X_A = 101$
 $Y_A = \alpha^{X_A} = \alpha^{101}$
 User B: $X_B = 42$
 $Y_B = \alpha^{X_B} = \alpha^{42}$

5. Encryption:
 User A: $M = \alpha^4$
 Public directory: $\alpha = x + 1, GF(2^8)$
 $Y_B = \alpha^{42}$
 User B: $X_B = 42$
 $Y_B = \alpha^{42}$ (aus 4.)

Encryption:
 $r = \alpha^{91} = \alpha^{2^3}$
 $C_2 = M * Y_B^r = \alpha^4 * (\alpha^{42})^{91} = \alpha^4 * \alpha^{3822 \text{ mod } 255} = \alpha^4 * \alpha^{252}$
 $= \alpha^{256 \text{ mod } 255} = \alpha$

6. Decryption:
 $Z^{-1} = r^{-X_B} = (\alpha^{91})^{-42} = \alpha^{-3822 \text{ mod } 255} = \alpha^3$
 $M = C_2 * Z^{-1} = \alpha * \alpha^3 = \alpha^4$

Solution 9-7 cont.:



Online ad-hoc Interactive Examples

Online Tutorial 1. 30.03.2021:

El-Gamal crypto system is set up, use the element $\alpha=35$ as a public element and compute the DH public keys Y_A and Y_B for users A and B having the secret keys $X_A=21$ and $X_B=29$.

1- Generate a prime number by using Pocklington Theorem:

$N = FR + 1$
 $N = 2^2(41) + 1 = 83, F = 41$ and $R = 2$

83 is a prime if all the following conditions hold :

Proof: select $a = 2$ ($1 < a < 83$) remark a is an integer prime to N

- $a^{N-1} \equiv 1 \pmod{N}$
 $2^{82} \equiv 1 \pmod{83}$ (or in Z_{83}) is true
- $\text{gcd}(a^{(N-1)/q_i} - 1, N) = 1$
 $\text{gcd}(2^{(83-1)/41} - 1, 83) = \text{gcd}(2^2 - 1, 83) = \text{gcd}(4, 83) = 1$ is true
- $F = 41 > \sqrt{83} \Rightarrow 41 > 9.11$ is true

As all conditions hold, $\Rightarrow p=83$ is for sure a prime

3- Compute and verify the Signature S_A according to ElGamal signature scheme for the same message $M=60$

User A signs M
 $X_A = 21$
 $Y_A = \alpha^{X_A} = 35^{21} \text{ mod } 83$
 $\alpha = 35$ primitive element in $GF(83)$
 $Y_B = \alpha^{X_B} = 52$
 $Y_C = \alpha^{X_C} = 80$
 Verifier
 $X_C = 29$
 $Y_C = \alpha^{X_C} = 35^{29} \text{ mod } 83$



Signed Message S_A

m	u	a1	a2	b1	b2	q	r	INVERSE VALUE = 82	GCD
82	9	1	0	0	1	9	1		
9	1	0	1	1	-9	9	0	INVERSE= -9	GCD= 1

2- Compute the ElGamal cryptogram CA using $GF(83)$ for the message $M=60$ sent from A to B using a random $R=9$.

